

# Algorithms for Genome Rearrangements

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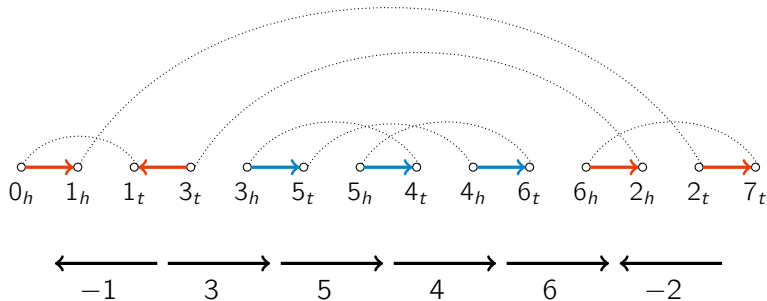
Lecture 4 – Sorting by Signed Reversals II

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# Quick Recap

BP Graph, **oriented** and **unoriented** components:



# Quick Recap

- Sorting is equivalent of increasing # of cycles in BP graph
  - In *oriented (good) components* – at least 1 oriented edge – this is always possible (safe reversals).
  - In *unoriented (bad) components*, we will need **extra operations**.
- If there are only oriented components in the BP graph:

$$d = N - C$$

- If there are also unoriented components:

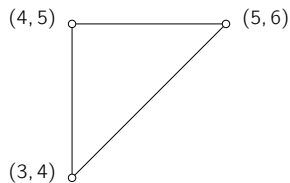
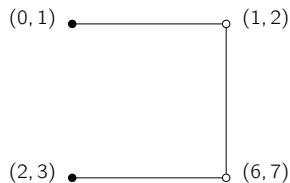
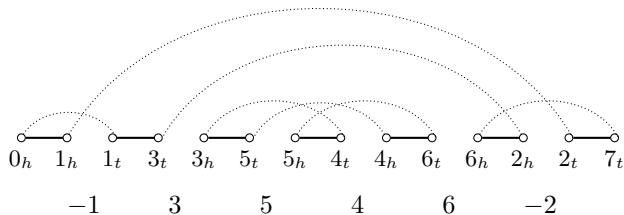
$$d = N - C + k,$$

where  $k$  is the minimum cost of these *extra* operations.

# Overlap Graph $O(\pi)$

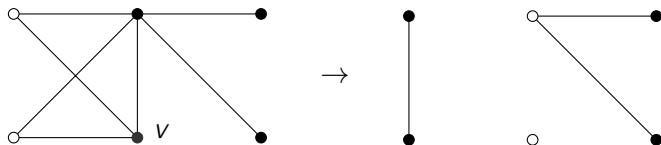
- Each vertex in  $O(\pi)$  corresponds to a gray edge in  $BP(\pi)$ .
  - Oriented edges  $\rightarrow$  Black vertices (odd degree).
  - Unoriented edges  $\rightarrow$  White vertices (even degree).
- If two edges in  $BP(\pi)$  overlap, there is an edge in  $O(\pi)$  between the corresponding vertices.

# Overlap Graph $O(\pi)$



# Effect of Reversal in the Overlap Graph

- 1 The subgraph induced by  $v$  and its neighbours is **complemented**.
- 2 All neighbours of  $v$  have their orientation inverted.

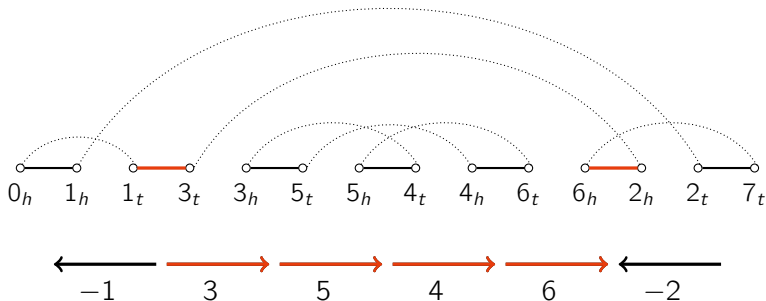


- The **score** of a reversal induced by  $v$  is

$$s(v) = T + U - O - 1$$

# Sorting Unoriented Components

- Let's analyse the effect that reversals have on cycles of  $BP(\pi)$ .
- Reversals change # of cycles by  $-1$ ,  $0$ , or  $+1$ .
- What happens exactly when we apply a reversal **defined** by two black edges?



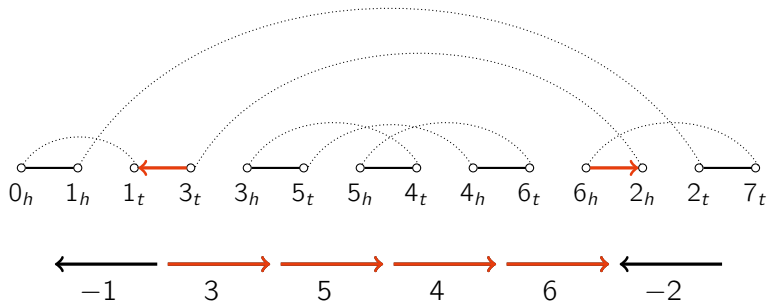
# Reversals and effect on cycles

- 1 Edges are on the **same cycle**:
  - **Type I**: Divergent edges: breaks the cycle.  $\Delta C = +1$ .
  - **Type II**: Convergent edges:  $\Delta C = 0$ , may change cycle orientation.
- 2 Edges on **different cycles**:
  - **Type III**: Merges the two cycles.  $\Delta C = -1$ .

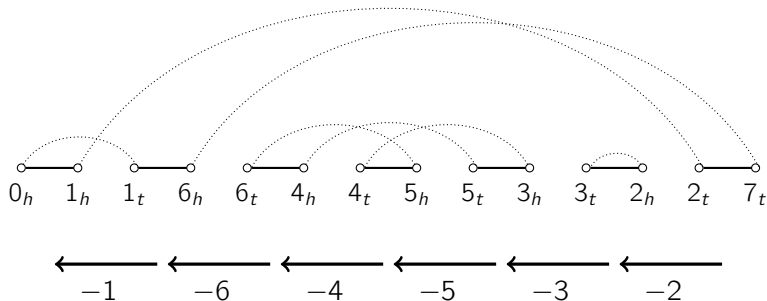
So far, we only used **Type I** operations, to sort oriented components.



# Type I - Same Cycle, divergent

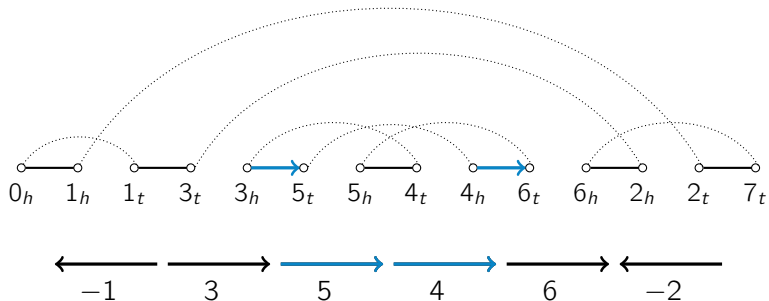


## Type I - Same Cycle, divergent

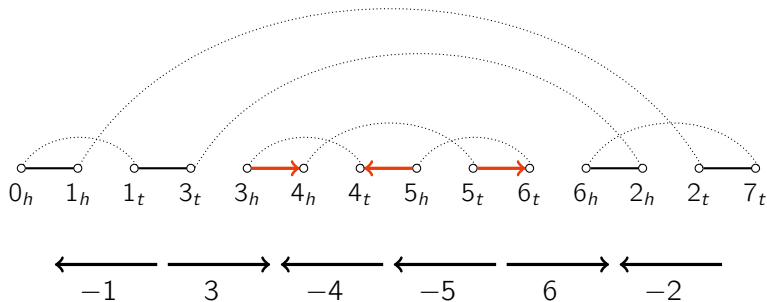


This reversal increases the number of cycles by one,  $\Delta C = +1$ .

## Type II - Same Cycle, convergent

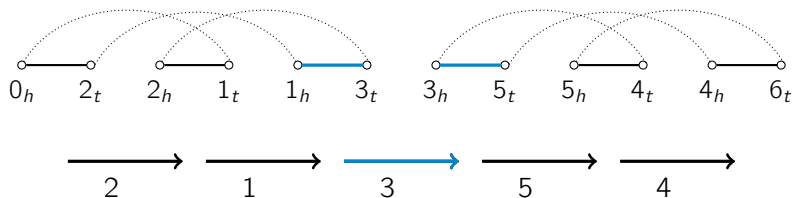


## Type II - Same Cycle, convergent

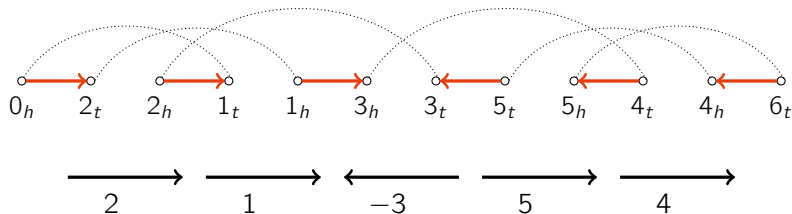


Does not change number of cycles ( $\Delta C = 0$ ), but the cycle is **oriented**.

# Type III - Different Cycles



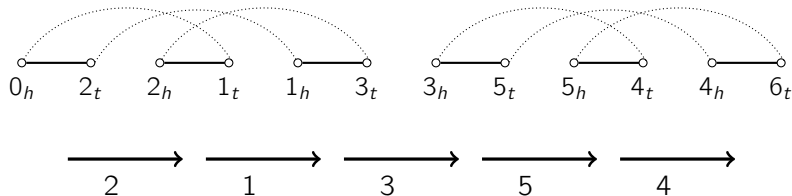
## Type III - Different Cycles



Merges the two cycles, decreasing the number of cycles by one ( $\Delta C = -1$ ), but the new cycle is **oriented**.

# Extra Operations

- How many extra operations do we need to sort unoriented components?



```
java InversionVisualisation L2/2unoriented.txt
```

## Extra Operations

- Applying one reversal in each cycle, orients both cycles, with 2 extra operations:

$$d = N - C + 2$$

- Applying one reversal merging both cycles, creates one new **oriented** cycle. Only one operation, but also one less cycle:

$$d = N - (C - 1) + 1 = N - C + 2$$

- In both cases, 2 extra operations. Does this mean that

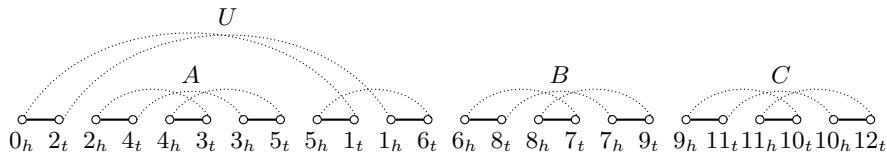
$$d = N - C + K$$

where  $K$  is the number of unoriented components? **Almost...**



# Definitions

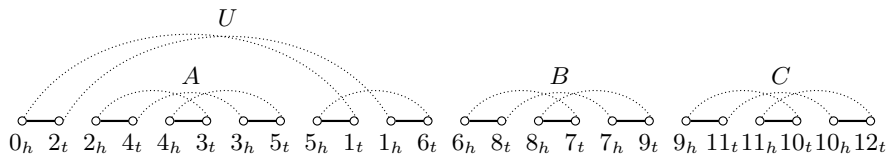
- A Component  $U$  **separates** two other components  $A$  and  $B$  if any edge from a vertex from  $A$  to  $B$  would cross an edge of  $U$ .



- $U$  separates  $A$  and  $B$ . (Also  $A$  and  $C$ ).

# Definitions

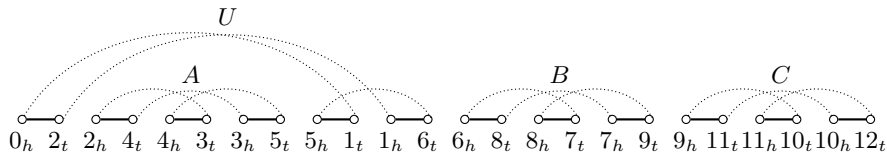
- A **hurdle** is an unoriented component that does **not** separate other two unoriented components.



- $A, B$  and  $C$  are hurdles.

# Definitions

- A **super-hurdle** is a hurdle that, when removed, causes the creation of a new hurdle.



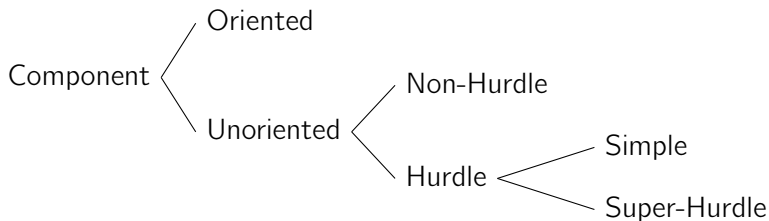
- $A$  is a super-hurdle.  $B$  and  $C$  are called *simple* hurdles.

- Why are these definitions important? Because except for one very rare special case, we have

$$d = N - C + H$$

where  $H$  is the number of hurdles.

# BP Graph – Component Types



# Reversal Types

- **Type I: Oriented Reversal:**  $\Delta C = +1$ .
  - Edges on same cycle, divergent.
- **Type II: Hurdle Cutting:**  $\Delta C = 0$ ,  $\Delta H = -1$ .
  - Edges on same cycle (hurdle), convergent.
- **Type III: Hurdle Merging:**  $\Delta C = -1$ ,  $\Delta H = -2$ .
  - Edges on different cycles (hurdles).

# Separating component

- Why is a separating component not a Hurdle?
- Because it can be oriented by a **Hurdle Merging** of two hurdles that is separates.

```
java InversionVisualisation L2/sep-hur-example.txt
```

# Super-hurdles: Problems might occur

- Cutting a super-hurdle is bad.
- Merging hurdles that are separated from a super-hurdle can cause the separating component to become a hurdle.

```
java InversionVisualisation L2/sep-hur-example.txt
```

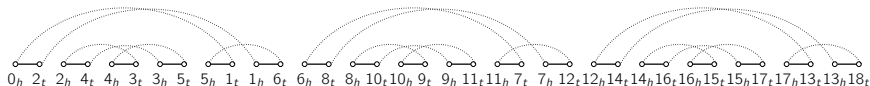


## Super-hurdles: Problems might occur

- How to avoid those problems?
- When there is an odd # of hurdles, cut a simple hurdle.
- When there is an even # of hurdles, merge opposite hurdles.
- Can we always do that? No... meet the **fortress**!

# Fortresses

- A **fortress** is a permutation that has an odd number of hurdles, and all are super-hurdles.



In this kind of permutation, there is no way to avoid an **extra operation**. A hurdle cut or hurdle merge will always create a new hurdle.

```
java InversionVisualisation L2/fortress.txt
```

# Reversal Distance - Complete equation

## Theorem (Reversal Distance, HP 95)

*The reversal distance of a permutation  $\pi$  is given by*

$$d(\pi) = N - C + H + F$$

*where:*

- *$N$  is the number of genes*
- *$C$  is the number of cycles in  $BP(\pi)$*
- *$H$  the number of hurdles in  $BP(\pi)$*
- $F = \begin{cases} 1, & \pi \text{ is a fortress} \\ 0, & \text{otherwise} \end{cases}$

# Reversal Distance - Complete Algorithm

```
1: procedure ReversalSort( $\pi$ )
2:   while  $\pi \neq$  identity do
3:     if  $\exists$  oriented component in  $BP(\pi)$  then
4:        $\rightarrow$  Apply a max score oriented reversal – Type I
5:     else if even # of hurdles then
6:        $\rightarrow$  Apply a Hurdle Merging on opposite hurdles – Type III
7:     else if  $\exists$  simple hurdle then
8:        $\rightarrow$  Apply a Hurdle Cutting on a simple hurdle – Type II
9:     else
10:       $\rightarrow$  Merge any two super hurdles (Fortress)
11:    end if
12:  end while
13: end procedure
```