## Algorithms in Comparative Genomics

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Exercise sheet 4, 07.11.2019

## Exercise 1 (Genome halving under DCJ)

Given an all-duplicates genome with three linear chromosomes

 $G = [3 \ 5 \ -4 \ 2 \ -5 ] [2 \ 1 ] [3 \ 4 \ 1 ],$ 

solve the genome halving problem under the DCJ distance, i.e., find a perfectly duplicated genome H with smallest DCJ distance to G.

## Exercise 2 (Guided genome halving under DCJ)

Research the guided genome halving problem under the DCJ distance and its complexity.

## Exercise 3 (Genome halving under the breakpoint distance)

The following paragraph about the genome halving problem for the breakpoint distance d is from the paper "Multichromosomal median and halving problems under different genomic distances" by Eric Tannier, Chunfang Zheng and David Sankoff, BMC Bioinformatics 10:120, 2009. (Notation: For a gene g, g1 and g2 denote the two paralogous copies of g; similarly, for a gene extremity x, the two copies are denoted x1 and x2.)

Let  $\Delta$  be an all-duplicates genome on a gene set  $\mathcal{G}$ , and G be the graph on the vertex set containing (1) all the extremities of the genes in  $\mathcal{G}$ , and (2) one supplementary vertex  $t_x$  for every gene extremity x. For any pair of gene extremities x, y, draw an edge in G weighted by zero, one or two according to the number of adjacencies in  $\Delta$  among x1y1, x1y2, x2y1, and x2y2. Now for any vertex x, draw an edge  $xt_x$  weighted by half the number of telomeres among x1 and x2 in  $\Delta$ . Finally, put an edge of weight 0 between  $t_xt_y$  for all pairs of gene extremities x, y.

For a genome M on  $\mathcal{G}$ , define a perfect matching, also called M, by including edges xy and  $t_x t_y$  for each adjacency xy, and an edge  $xt_x$  for each telomere x. Let w(M) be the weight of the matching M.

1. Prove the following:

Claim. For a genome M on  $\mathcal{G}$ , the perfect matching M thus constructed satisfies  $w(M) = 2n - d(\Delta, M)$ .

2. How can this construction be used to solve the genome halving problem under the breakpoint distance?

*Hint:* Have a closer look at the all-duplicates genome  $\Delta = \begin{bmatrix} 1 & 2 & 3 \end{bmatrix} \begin{bmatrix} 1 & -3 & -2 \end{bmatrix}$ .